



Global Optimization of ODE constrained network problems

Marc Pfetsch

Joint work with: Oliver Habeck and Stefan Ulbrich

TU Darmstadt



Discrete
Optimization

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Optimization using the Example of Gas Networks

- 1 General Theory
 - Abstract Model and Relaxations
 - Solving Strategies
- 2 Application – Stationary Gas Transport
 - Model
 - Relaxation of the ODE
- 3 Implementation – for Stationary Gas Transport
 - The Algorithm
 - Results

We consider the ODE constrained optimization problem:

$$\begin{aligned} \min \quad & C(x, y^0, y^S, z) \\ \text{s.t.} \quad & G(x, y^0, y^S, z) \leq 0, \\ & \partial_s y(s) = f(s, x, y(s)), & s \in [0, S] & \quad (\mathcal{P}_{ode}) \\ & y^0 = y(0), \quad y^S = y(S), \\ & x \in X, \quad y^0 \in Y^0, \quad y^S \in Y^S, \quad z \in Z, \end{aligned}$$

- ▷ $Y^0, Y^S \subset \mathbb{R}^n$ and $X \subset \mathbb{R}^k$ are polytopes and $Z \subset \mathbb{Z}^m$ is bounded.
- ▷ $C: X \times Y^0 \times Y^S \times Z \rightarrow \mathbb{R}$ and $G: X \times Y^0 \times Y^S \times Z \rightarrow \mathbb{R}^l$ are continuously differentiable and possibly nonlinear.
- ▷ Values of y only need to be known for finitely many points $y(0), y(S) \in \mathbb{R}^n$.
- ▷ Examples: Stationary networks (gas, water, (electricity), ...)



- ▶ Standard approach is “Discretize then optimize”:
discretize ODE (add N variables), use numerical approximation scheme and obtain MINLP. Solve it and possibly refine.
- ▶ Our goal: compute globally optimal solution without explicitly discretizing.
- ▶ Workhorse: good convex/concave under/over-estimators.
- ▶ Will show: Easy to construct in certain cases, e.g., gas optimization.
- ▶ Yields globally convergent solution algorithm.

Equivalent reformulation with the solution mapping $F: (x, y^0) \mapsto y^S$:

$$\begin{aligned} \min \quad & C(x, y^0, y^S, z) \\ \text{s.t.} \quad & G(x, y^0, y^S, z) \leq 0, \\ & y^S - F(x, y^0) = 0, \\ & x \in X, y^0 \in Y^0, y^S \in Y^S, z \in Z. \end{aligned} \tag{P}$$

Nonconvex Relaxation of (\mathcal{P})

Assume that there exist $F^\ell : X \times Y^0 \times \mathbb{N}^n \rightarrow \mathbb{R}^n$ and $F^u : X \times Y^0 \times \mathbb{N}^n \rightarrow \mathbb{R}^n$, with

$$F^\ell(x, y^0, N) \leq F(x, y^0) \leq F^u(x, y^0, N)$$

for all $x \in X$ and $y^0 \in Y^0$. Furthermore, on the polytopes X , Y^0 the functions F^ℓ and F^u converge uniformly to F for $N \rightarrow \infty$, $i = 1, \dots, n$.

This yields the relaxation

$$\begin{aligned} \min \quad & C(x, y^0, y^S, z) \\ \text{s.t.} \quad & G(x, y^0, y^S, z) \leq 0, \\ & F^\ell(x, y^0, N) \leq y^S \leq F^u(x, y^0, N), \\ & x \in X, y^0 \in Y^0, y^S \in Y^S, z \in Z. \end{aligned} \tag{\mathcal{P}_r(N)}$$

Assume that there exist

- ▷ convex underestimators $\check{C} \leq C$, $\check{G} \leq G$, and $\check{F}^\ell \leq F^\ell$
- ▷ and a concave overestimator $F^u \leq \hat{F}^u$.

Hence,

$$\begin{aligned} \min \quad & \check{C}(x, y^0, y^S, z) \\ \text{s.t.} \quad & \check{G}(x, y^0, y^S, z) \leq 0, \\ & \check{F}^\ell(x, y^0, N) \leq y^S \leq \hat{F}^u(x, y^0, N), \\ & x \in X, y^0 \in Y^0, y^S \in Y^S, z \in \text{conv}(Z) \end{aligned} \tag{\mathcal{P}_{cv}(N)}$$

is a convex relaxation of $(\mathcal{P}_r(N))$.

↪ we can solve $(\mathcal{P}_r(N))$ with spatial branch-and-bound.



Definition

A vector $(x, y^0, y^S, z) \in X \times Y^0 \times Y^S \times Z$ is (ϵ, δ) -optimal, if

- ▶ each constraint is violated by a maximum of δ
- ▶ and the objective function satisfies $C(x, y^0, y^S, z) \leq C^* + \epsilon$.

Goal: Find $(\epsilon, \delta_1 + \delta_2)$ -optimal solution of (\mathcal{P}) .

- ▶ Start spatial branch-and-bound with some initial $N \in \mathbb{N}^n$ and feasibility tolerance δ_1 .
- ▶ Solve the convex relaxation $(\mathcal{P}_{cv}(N))$ in a node.

- ▶ Check if

$$\|F^u(x, y^0, N) - F^\ell(x, y^0, N)\|_\infty \leq \delta_2$$

holds for the current solution of the relaxation.

- ▶ Increase N if necessary, improve the under- and overestimators \check{F}^ℓ , \hat{F}^u , and solve the convex relaxation again.
- ▶ Continue like in the “normal” spatial branch-and-bound algorithm.

Theorem

For a nested sequence of nodes $\mathcal{F}_k = X_k \times Y_k^0 \times Y_k^T \times Z_k$ with $\mathcal{F}_{k+1} \subseteq \mathcal{F}_k$ let

- ▷ $\lim_{k \rightarrow \infty} \text{diam } \mathcal{F}_k = 0$
- ▷ and the estimators become exact in the limit, i.e., the estimators converge to the original functions on \mathcal{F}_k for $k \rightarrow \infty$.

Then the solving strategy produces an $(\epsilon, \delta_1 + \delta_2)$ feasible point of (\mathcal{P}) or shows that it is infeasible in finite time.

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- ▷ Gas networks are given by directed graphs $\mathcal{G} = (\mathcal{V}, \mathcal{A})$.
- ▷ The nodes are junctions of the network.
- ▷ The arcs are
 - (control)valves: linear constraints with binary variables
 - resistors: (non)linear constraints with indicator variables
 - compressors: (here) polyhedral approximation of operating range
 - pipes: ordinary differential equation

Stationary isothermal Euler equation with no inclination and $\frac{|v|}{c} = \frac{c|q|}{Ap} \leq 0.8$:

$$\partial_x p(x) = -\frac{\lambda c^2 q|q| p(x)}{2D (A^2 p(x)^2 - c^2 q^2)} =: \phi(p(x), q).$$

Variables:

- ▷ pressure $p(x)$
- ▷ mass flow q

Constants:

- ▷ friction coefficient λ
- ▷ speed of sound c
- ▷ diameter D
- ▷ cross sectional area A
- ▷ length of the pipe L

For $q \geq 0$ the solution $p(x)$ is concave and non increasing.

Thus, we want to solve the following optimization problem:

$$\begin{aligned} \min \quad & C(p, q, z) \\ \text{s.t.} \quad & G(p, q, z) \leq 0, \\ & \partial_x p_a(x) = \phi_a(p_a(x), q_a) \quad \forall a \in \mathcal{A}_{\text{pipe}} \subseteq \mathcal{A}, \\ & p_u = p_a(0), \quad p_v = p_a(L_a) \quad \forall a = (u, v) \in \mathcal{A}_{\text{pipe}}, \\ & \underline{q}_a \leq q_a \leq \bar{q}_a \quad \forall a \in \mathcal{A}, \\ & \underline{p}_u \leq p_u \leq \bar{p}_u \quad \forall u \in \mathcal{V}, \\ & z \in Z \subset \mathbb{Z}^m \end{aligned}$$

with spatial branch-and-bound.



Theorem

Given a discretization $0 = x_N < x_{N-1} < \dots < x_0 = L$ with $h_{i+1} = x_i - x_{i+1}$ for $i = 0, 1, \dots, N - 1$. If $q \geq 0$, bounds on $p(0)$ are given by the following methods:

▷ *Upper bounds through the trapezoidal rule:*

$$p_0^u = p_{out}, \quad p_{i+1}^u + \frac{1}{2} h_{i+1} \phi(p_{i+1}^u, q) = p_i^u - \frac{1}{2} h_{i+1} \phi(p_i^u, q), \quad \forall i = 0, 1, \dots, N - 1.$$

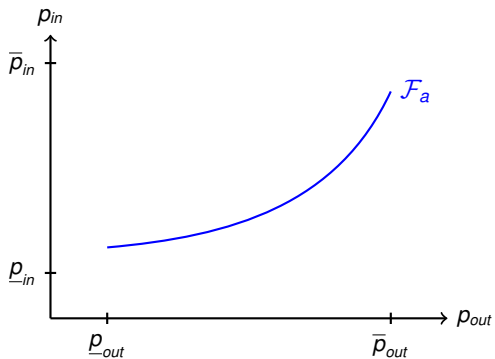
▷ *Lower bounds through the explicit midpoint method:*

$$p_0^l = p_{out}, \quad p_{i+1}^l = p_i^l - h_{i+1} \phi(p_i^l - \frac{1}{2} h_{i+1} \phi(p_i^l, q), q), \quad \forall i = 0, 1, \dots, N - 1.$$

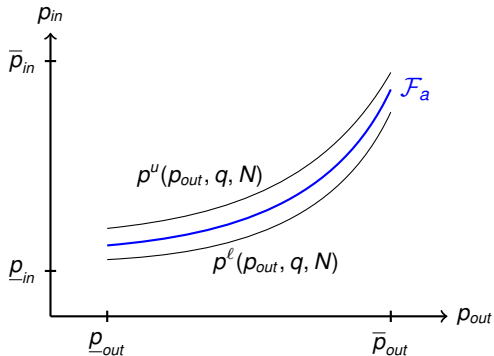
↪ yields (nonconvex) relaxation of \mathcal{F}_a .

Critical property: **nonnegative local truncation error.**

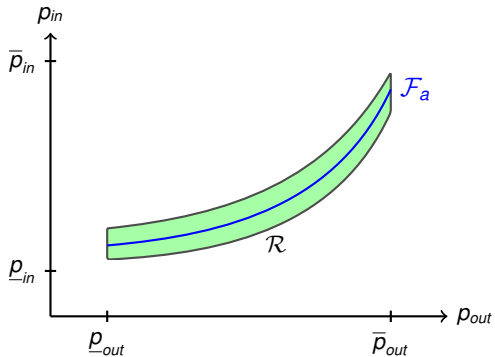
Illustration



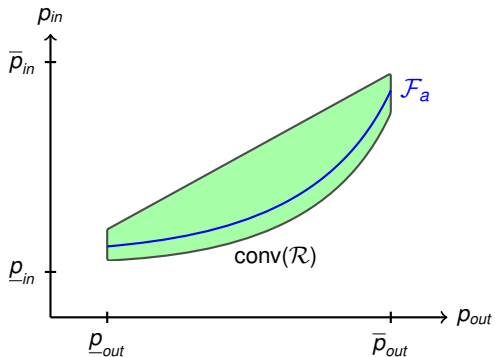
Illustration



Illustration



Illustration





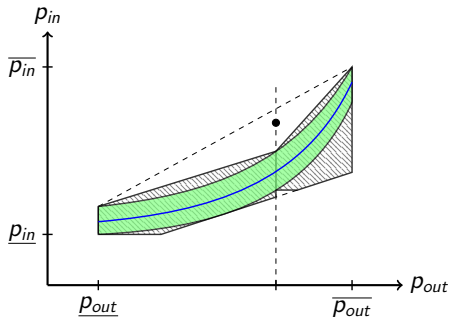
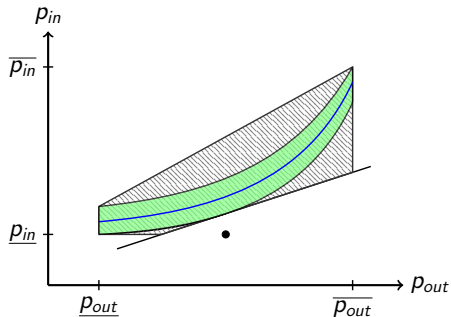
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- ▷ We use the branch-and-bound framework SCIP.
- ▷ We implemented a constraint handler for the ODEs.
- ▷ Test instances are available at <http://gaslib.zib.de/>.

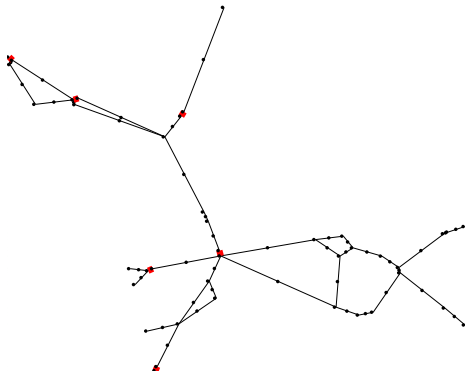
We perform the following steps in the nodes of the branch-and-bound tree:

1. Perform bound propagation (use estimators)
2. Solve LP-relaxation
3. Enforce feasibility of “most infeasible” pipe by one of the following steps:
 - 3.1 Fixing the direction of the flow by branching.
 - 3.2 Adding (parts of) the concave overestimator.
 - 3.3 Branching w.r.t. input pressure, output pressure or massflow.
 - 3.4 Adding a gradient cut (convex underestimator).
4. Resolve new LP-relaxation or choose a new node.

Illustration



Example: GasLib-40



- ▷ 40 nodes
- ▷ 39 pipes
- ▷ 6 compressors
- ▷ 12 binary variables for compressors

Results

objective: minimize	$-\sum_{v \in V} p_v$	$\sum_{v \in V} q_v^\pm p_v$	$\sum_{a \in \mathcal{A}_{CS}} z_a$
solving time (seconds)	877.74	542.48	65.08
processed Nodes	12,444	2790	263
branchings on flow (%)	46.2	81.0	84.2
branchings on pressure (%)	52.9	0.0	8.8
branchings on binary variables (%)	0.9	19.0	7.0
leaves	5900	1412	73
cut offs by propagation	5297	1149	38
bound changes by propagation	32,444	63,317	6051
added overestimators	7362	2385	985
added underestimators	6360	4281	488

$$\delta_1 = \epsilon = 10^{-6}, \delta_2 = 10^{-4}$$

Future theoretical goals:

- ▷ Extend to case with inclination.
- ▷ Include temperature?
- ▷ Transient case (much more difficult).

Future implementation issues:

- ▷ Better branching rules
- ▷ Additional primal heuristics
- ▷ Avoid cycle flows
- ▷ Cuts on binary variables
- ▷ Larger instances



O. Habeck, S. Ulbrich, and M. E. Pfetsch.

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Preprint, TRR 154, 2017.

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